Propositional logic and set theory pdf

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Set theory is a branch of mathematical logic. Therefore, it is natural that logical language and symbols are used to describe sets. In this section, we will look at the basic logical symbols and ways of defining sets. Propositional Logic A proposition is a declarative statement which is either true or false. If a proposition is true, then we say it has a truth value of true. Respectively, if a proposition is false, its truth value is false. So, for example, the following statements have a truth value of true: The Earth revolves around the Sun; (10 + 3 = 13;) If (x) is an even integer, then  $({x^2})$  is also even. Examples of false propositions: An electron is heavier than a proton; (1 + 2 gt 3;) (6) is a prime number. Not all sentences are propositions: \(x \gt 5\) (This may be true or false depending on \(x\)) Is it raining? (This is a question, not a declarative sentence) Mondrian paintings are too abstract?) To represent propositions, we denote them by letters. The most common letters are \(p,\) \(r,\) \(s,\) \(t.\) Using logical operators or connectives, we can build compound propositions. Logical Operators and Truth Tables Let \(p\) and \(q\) be two propositions. Each of these statements can take two values - true (\(T\)) and false (\(F\)). So there are \(4\) pairs of input values: \(TT,\) \(TF,\) and \(FT,\) (q.) The truth values of the proposition (r) can take different values. Figure 1. There are total  $(2^4 = 16)$  possible output combinations (truth functions) for (2) binary input variables. Each of these combinations is represented by a certain logical operator. Further we'll look at the most important operators. Negation Negation is a unary logical operator. If \(p\) is a proposition, then the negation of \(p\) is called not \(p\) and is denoted by \(\lnot p.\) To represent the meaning of a logical expression, it is convenient to use a truth table. proposition(s). In case of the negation operator, the truth table is very simple: As you can see, the logical negation operator reverses the truth value of the input proposition. Example 1: \(p:\) A trapezoid is not a quadrilateral (false) Example 2: \(p:\) \(2\) is a prime number (true) \((\lnot p:\) A trapezoid is not a quadrilateral (true) \((\lnot p:\) A trapezoid is not a quadrilateral (true) \((\lnot p:\) A trapezoid is not a quadrilateral (false) Example 2: \(p:\) A trapezoid is a quadrilateral (true) \((\lnot p:\) A trapezoid is not a quadrilateral (true) \((\lnot p:\) A trap prime number (false) Consider now some binary operators. Conjunction If \(p\) and \(q\) are true only when both \(p\) and is denoted by \(p \land q.) is true only when both \(p\) and \(q\) are true. Otherwise, it is false. Thus, the truth table for conjunction looks as follows: Example 1: \(p:\) United Kingdom is a member of European Union (false) \(q:\) \(3\) is prime (true) \(q:\) \(3\) is prime and odd (true) \(p\land q:\) United Kingdom and Ireland are members of European Union (false) A (q:\) Ireland is a member of European Union (false) (q:\) (3\) is odd (true) \(p\land q:\) (3\) is odd (true) \(p\land q:\) United Kingdom and Ireland are members of European Union (false) (q:\) (3\) is prime (true) \(q:\) (3\) is prime (true) (q:\) (3\ disjunction means \(p\) or \(q\) and is denoted by \(p \lor q.) The disjunction \(p \lor q\) is true when either \(p\) is true, or both are true. It is false, if both \(p\) and \(q\) are false. Truth table for disjunction: Example 1: \(p:\) A proton has a negative charge (false) \(q:\) A neutron has a negative charge (false) \(p \lor q:\) A proton or a neutron has a negative charge (false) Example 2:  $(p) ((\frac{2}{3}))$  or  $((\frac{2}{3}))$  or  $((\frac{2}{3})$ denoted by \(p \to q.) This logical operator is also called just a conditional \(p \to q\) is true, then the conditional \(p \to q\) is false, then the conditional \(p \to q) is assumed to be true by default. Here is the truth table for conditional \(p \to q) is assumed to be true by default. sentences, some of them are listed below: (p) (q) (q) is a sufficient condition for (q) (q) (q) is a necessary condition for (p) (q) (q) is divisible by (2) and (3) (q) (q) is divisible by (2) (q) (q) (q) is divisible by (2) (q) (q) (q) (q) (q) (q)is true, and  $(p \ q)$  is true. (x = 13:) (q) is the proposition  $(q \ q)$  is true, then the contrapositive is also true. If the converse is true, then the inverse is also true. The special conditional operators are defined by the following truth table: Example: Statement \(p\) A quadrilateral is a rectangle (false) Statement \(p\) degrees (true). Conditional statement \(p\) If a quadrilateral is a rectangle, then the sum of its interior angles is \(360\) degrees (true). Converse of \(p \to q:\) \(q \to p\) If the sum of interior angles of a quadrilateral is a not a rectangle (false). Inverse of \(p \to q:\) \(eg p \to eg q\) If a quadrilateral is a not a rectangle, then the sum of its interior angles is not \(360\) degrees (false). Contrapositive of \(p \to q:\) \(eg q \to eg p\) If the sum of interior angles of a quadrilateral is not \(360\) degrees, then it is not a rectangle (true). Material Biconditional of \(p\) and is denoted by \(p \leftrightarrow q.\) The biconditional statement has the same truth value as the compound logical operator  $(\left\{q \ p \ p \ right\right), \$  which means (p) implies (q) and (q) implies (p,) The scalar triple product of three vectors is zero (true)  $(p \ right), \$ Three vectors are coplanar if and only if their scalar triple product is zero (true) (p\) The tennis match will be played outdoors (false) \(q:) It is raining (false) Logical Equivalence Propositions \(p) and \(q) are said to be logically equivalent if they have the same truth tables. The logical equivalence of \(p\) and \(q\) is denoted as \(p \equiv q,\) or sometimes by \(\Leftrightarrow\) depending on the notation being used. Predicates and Quantifiers So far, we considered propositional logic, we introduce the concept of predicate. A predicate is a logical statement that contains one or more variables or parameters. The predicates are denoted by a capital letter and the variables or parameters. The predicate is a logical statement that contains one or more variables or parameters. The predicate is a logical statement that contains one or more variables or parameters. The predicate depends on the values of its variables. Example 1: Predicate \  $(P\left(\frac{x,y}\right) (Q\left(\frac{x,y}{right})\right) (Q\left(\frac{x,y}{right}\right)) (Q\left(\frac{x,y}{righ$  $\left(\frac{1^2} + \frac{1^2} +$ complex predicates. There is also an additional operation defined on predicates and called quantification. Quantification allows us to specify the extent of validity of a predicate should hold. There are two types of quantifiers in predicate logic - universal quantifier and existential quantifier. Universal Quantifier The universal quantifier is used to express sentences with words like all or every. It is denoted by the symbol (forall x P(t(x)right)) is true". The domain is called the universe of discourse or domain of discourse. Existential Quantifier The existential quantifier is used to express sentences with words like some or there is a. It is denoted by the symbol \(\exists.) The notation \(\exists.) The n describing mathematical objects and for modeling the real world. Set Builder Notation The set builder notation is used to specify a set of objects by means of a predicate. The common notation includes \(3\) parts: a variable \(x,\) a colon or vertical bar separator, and a logical predicate \(P\left({x}\right):\) \[S = \left({ x|P\left(x \right)}) \[S = \left(x \right) \[S = \left(x \right)) \[S = \left(x \righ member of \(S\) \(x otin S\) \(x\) is not an element of \(S\) \(S\) is a subset of \(T\) \(S = T\) is equivalent to \(\left({S \subseteq T}\right)) \(S\) is a proper subset of \(T.) This means \(\left({S \subseteq T}\right)) \(S\) is a subset of \(T.) \(S) is a proper subset of \(T.) This means \(\left({S \subseteq T}\right)) \(S \subseteq T) \(S) is a subset of \(T.) \(S \subseteq T) \(S) is a proper subset of \(T.) This means \(\left({S \subseteq T}\right)) \(S \subseteq T) \(S) is a proper subset of \(T.) This means \(S) \(S \subseteq T) \(S) is a proper subset of \(T.) \(S \subseteq T) \(S) is a proper subset of \(T.) \(S \subseteq T) \(S) is a proper subset of \(T.) \(S \subseteq T) \(S) is a proper subset of \(T.) \(S \subseteq T) \(S) is a proper subset of \(T.) \(S \subseteq T) \(S \subset T) \( Page 2. Page 2 Click or tap a problem to see the solution. Prove De Morgan's Laws: \[eg \left( {p \lor q} \right) \left( {eg p} \right) \right) \right) \right) \left( {eg p} \right) \right) \right) \right) \right) \right) \right) \left( {eg p} \right) \right) \right) \right) \right) \right) \right) \rig goes to the gym \(4\) times a week. Nicole goes to the gym 2 times a week. Any sometimes goes to the gym 2 times a week, but not every week. Any sometimes goes to the gym 15 times a week. Nicole goes to the gym ((3\) is a math course, and \(M\left( {x,y} \right)) is the predicate "\(x\) will take \(y\)". Translate the following sentences into predicate notation: Every student will take Calculus. Every student will take Calculus. Every student will take a math course that everybody will take. There is a course that nobody will take. There is a course that nobody will take a math courses except Topology. Describe the following sets using set builder notation:  $(\left\{g, 27, -8, -1, 0, 1, 8, 27\} \right) (\left\{\frac{2}{3}, \frac{1}{2}, \frac{$ {eg p} \right) \land \left( {eg q} \right).] Solution. We verify the \(1\text{st}) De Morgan's Law using truth table. First we write all possible combinations of \(p\) and \(q\) The negations of \(p\) and \(q\) The negation of conjunction of \(p\) and \(q\) are their opposite values. \(left({p \land q}\right)). The expression \(\left({eg p} \right)) is disjunction of \(eg p\) and \(eg q.) 1st De Morgan's Law. We see that the last two columns of the truth table are the same. This proves the \(1\text{st}) De Morgan's Law. Similarly, we can verify the \(2\text{nd})) De Morgan's Law: 2nd De Morgan's Law. Show that the statement  $|s = \left( \left\{ p \right\} \right) |c q| right| |c q q| right| lor q q right| lor left( q q right| lor q right| lor q q right| lor q q right| lor q right| lor q q right| lor q right| lo$ disjunction of (p) and (eg q.) Then we find conjunction  $(\{eg p \mid and q\}.)$  The final statement (s) is disjunction of two previous propositions. A tautology statement  $(s = |eft( \{p \mid o q q\} \mid and q\} \mid s a tautology.$  Show that the statement  $(r = |eft( \{p \mid o q q\} \mid and q\} \mid s a tautology.$ (eg p \land q) \land (p \leftrightarrow q)\] is a contradiction. Solution. A contradiction is a statement that is always false. Let's construct a truth table to make sure that the statement \(r\) has only false values. The first expression \(eg p \land q) is conjunction of \(eg and \(q\). The final operation is conjunction of two previous propositions. A contradiction statement. The last column contains only false values. Hence, the proposition \(r = (eg p \land r) \right) \left( {p \lor q} \right) \land (p \left( {p \lor r} \right).] Solution. We construct a truth table to make sure that the left hand and right hand sides of the statement are equivalent. Let's denote \[s1 = p \lor \left( {p \lor r} \right),\;;s2 = \left( {p \lor r} \right) \] and find the truth values of \(s1\) and \(s2.\) Applying conjunction and disjunction operations we get the following table: A distributive law. The last two columns of the table are identical. This means that \[s1 \equiv s2,\;\; \Rightarrow p \lor \left( {p \lor r} \right)\] Let \(G\left( {x,y} \right)\) be the predicate notation: Nick goes to the gym \(y\) times a week." Write the following propositions in predicate notation: Nick goes to the gym \(y\) times a week." Write the following propositions in predicate notation: Nick goes to the gym \(y\) times a week." Write the following propositions in predicate notation: Nick goes to the gym \(y\) times a week." Write the following propositions in predicate notation: Nick goes to the gym \(y\) times a week." the gym \(4\) times a week. Nicole goes to the gym \(2\) or \(3\) times a week. Max goes to the gym 2 times a week. Max goes to the gym 2 times a week. Solution. Nick goes to the gym \(4\) times a week. Max goes to the gym 2 times a week. Max goes to the gym 15 times a week. Max goes to the gym 15 times a week. Max goes to the gym 2 times a week. Max goes to the gym 16 times a week. Max goes to the gym 15 times a week. Max goes to the gym 15 times a week. Max goes to the gym 15 times a week. Max goes to the gym 16 times a week. gym \(2\) or \(3\) times a week. \[G\left( {Nicole,2} \right) \lor G\left( {Nicole,3} \right)] Max goes to the gym 2 times a week, but not every week. \[G\left( {Max,0} \right)] Amy sometimes goes to the gym. \[\exists y G\left( {Amy,y} \right)] Some people go to the gym every day. \[\exists x G\left( {x,7} \right)] Hardly anyone goes to the gym 15 times a week. \[eg \exists x G\left( {x,15} \right)\] Suppose \(x\) is a student, \(y\)". Translate the following sentences into predicate notation: Every student will take Calculus. Every student will take a math course. Lora will take Calculus and Linear Algebra. There is a course that everybody will take. Tom will take a math course that nobody will take. Tom will take a math course sexcept Topology. Solution. Every student will take a math course sexcept Topology. Solution. Every student will take a math course that nobody will take. There is a course that nobody will take a math course sexcept Topology. Solution. Every student will take a math course sexcept Topology. [M\left( {Lora,Calculus} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that everybody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM\left( {x,y} \right)] There is a course that nobody will take. \[\exists y\forall xM Topology \right)\] Describe the following sets using set builder notation: (\left\{ { - 27, - 8, - 1,0,1,8,27} \right\}) ((\left\{ { - 27, - 8, - 1,0,1,8,27} \right\}) Solution. The set (\left\{ { - 27, - 8, - 1,0,1,8,27} \right\}) contains integers from (x = -3) to (x = 3) cubed.  $\left(\left\{ x \in \mathbb{Z}, \left(x^3\right) \in x \in \mathbb{Z}, \left(x^3\right) \in$  $1 \le \frac{1}{\pi \in 1}$  is a portion of the Fibonacci sequence (F n) from (n = 6) to  $(n = 11.) \left[ \frac{x \in 1}{x \in 1} \right]$  is a portion of the Fibonacci sequence (F n) from (n = 6) to  $(n = 11.) \left[ \frac{x \in 1}{x \in 1} \right]$ 

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